## **Pushdown Automata - Exercise**

Problem:

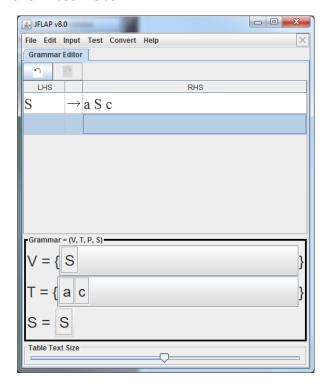
## Solution:

First, we examine the kinds of words produced by this set. One way to do that is to tabulate the different values of n, m, and k and list the corresponding words for the different values.

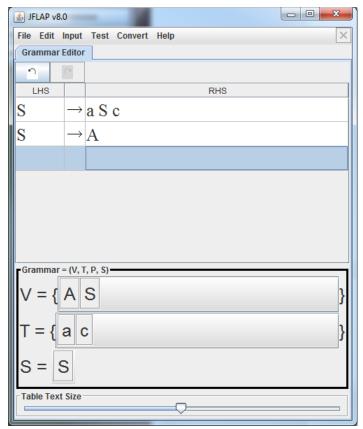
n	m	k	W
0	0	0	λ
0	1	2	bcc
1	0	1	ac
1	1	3	abccc
1	2	5	abbccccc

We see that for each a in the word, there is a c produced. For each b in the word, there are two b's. Also, a's, b's, and c's appear in that sequence; i.e., cccba is not accepted.

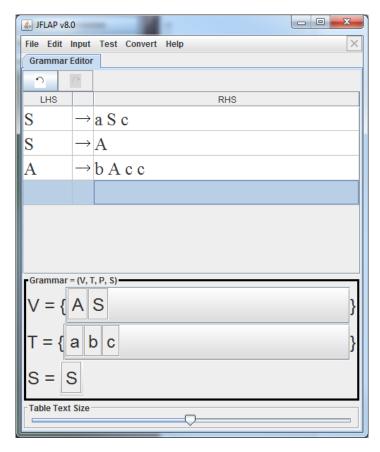
Open JFLAP and click on *Grammar*. Notice that the resulting window has a LHS column and a RHS column. Each row refers to one rule in the context-free grammar to be created. Using S as the start symbol, fill the first row with  $S \rightarrow$  aSc like so:



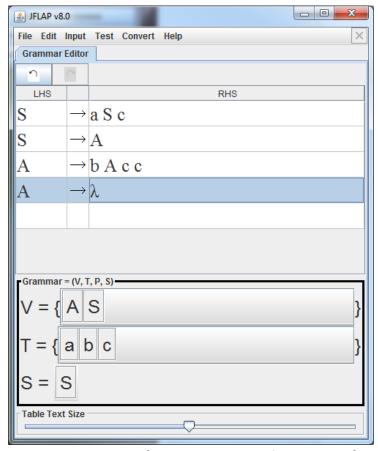
The S in the RHS allows the rule to become recursive. This rule states that for every a created, there is a c at the end of the word created. Next, we introduce a production rule to allow b's to be created. Note that once b's are being produced, we cannot create any more a's. We do this by adding a second rule,  $S \rightarrow A$ .



Next, we work on letting b's be created making sure that there are two c's for every b, A  $\rightarrow$  bAcc. This rule is also recursive allowing an infinite number of b's to be created.



Last, we add a rule to terminate the derivation,  $A \rightarrow \lambda$ .



That completes the grammar. Next, we test for grammar type with *Test > Test for Grammar Type*. JFLAP reports that it is a context-free grammar.

Lastly, we may test an input string. Click *Input > Brute Force Parse* and enter *abbccccc* as the input. Click *Complete*. JFLAP reports this string as *accepted*. The brute parse table also shows the 5 derivation levels and the progression from S to the test input.

